

ZK-STARK Theory & Implementation

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November 2021



Overview

- 1. My story and "red pill" moment
- 2. The Cambrian Explosion of ZKPs
- 3. ZK-STARKs unleashed
- 4. How to build a STARK?
- 5. [Fast RS IOPPs (FRI)] time permitting



My Story



My Research and Blockchain

- 2001: Postdoc at Harvard+MIT, Madhu Sudan suggested studying PCP length
- 2003-5: Short PCPs with poly-log query complexity [BS04]
 - Theoretical result, no practical application in sight
- 2008: Students start implementing it in code
 - Why? No clear reason
- 2009: Huge ERC funding (1.7M Euro), more implementation
 - Why? Still no good reason
- 2013: Bitcoin San Jose Conference
 - Red pill swallowed
 - Why?



Post Red Pill

- 2014: Zerocash academic paper
- 2015: Zcash launched
- 2013-16: Startup failed attempt
- 2018: Math breakthroughs, not well-received
 - FRI: Rejected from 3 conferences (including STOC/FOCS and ITCS, accepted to ICALP)
 - STARK: Rejected from 4 conferences (including CRYPTO, CCS, accepted to CRYPTO)
 - PCP Security: Rej from 3 conferences (gave up)



Meanwhile in Blockchain world...

- Zcash=> ZKP/ ZK-SNARKs hype
- Huge enthusiasm for ZK-STARKs
- 2018: StarkWare Founded
 - My co-founders: Alessandro Chiesa, Uri Kololdny, Michael Riabzev

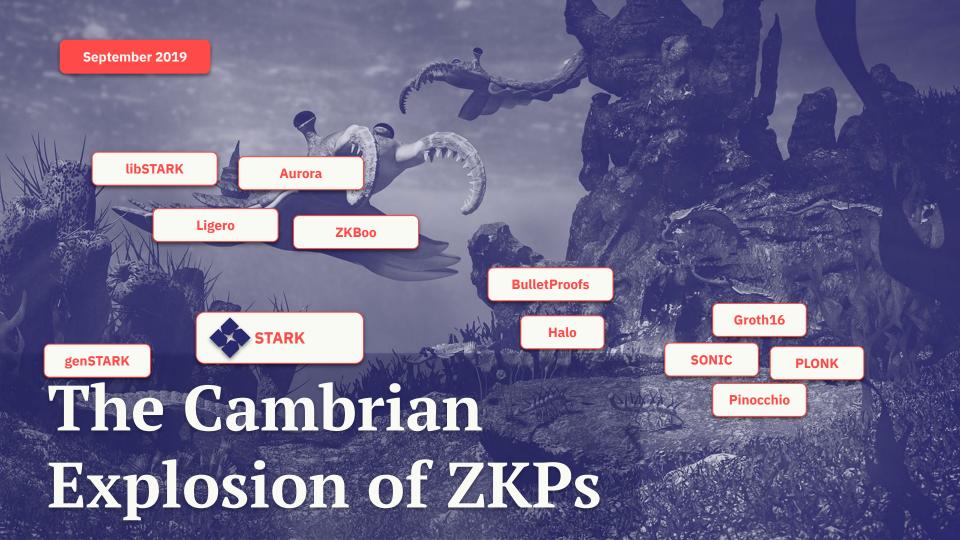
STARKWARE

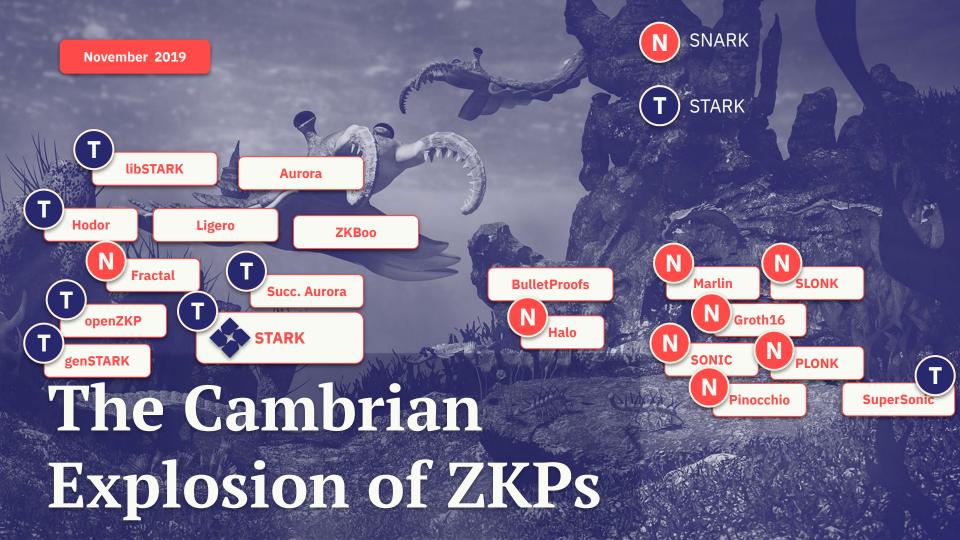
- \circ \$6M funding, followed by \$25M, ...
- At launch, still missing:
 - key math results: DEEP FRI, tight soundness analysis, ...
 - Accessibility: Cairo language, system, business model, product ...
 - But we knew very well what we'll do



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Proofs of Computational Integrity (CI)



Privacy (Zero Knowledge, ZK) Prover's private inputs are shielded



Scalability Exponentially small verifier running time* Nearly linear prover running time*



Universality Applicability to general computation



Transparency No toxic waste (i.e. no trusted setup)



STARK

*With respect to size of computation



Proofs of Computational Integrity (CI)



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(ZK)-STARK

*With respect to size of computation



STARK vs. SNARK - emphasizing different aspects



STARKs must be

Transparent no trusted setup

Scalable: logarithmic verifying time **and** nearly-linear proving time

Succinct setup, at most logarithmic time

Non-interactive STARKs are SNARKs (transparent ones)

Transparent SNARKs w/ succinct setup are STARKs



SNARKs must be

Noninteractive: pf is single message (after preprocessing)

Succinct: logarithmic verifying time

Setup can take linear time (and more)



1) Arithmetization

Arithmetization Converts ("reduces") Computational Integrity problems to problems about local relations between a bunch of polynomials

Example: For public 256-bit string *z*, Bob claims knows a SHA2-preimage of *z*

Pre-arithmetization claim	Reduction	Post-arithmetization claim	Theorem
"I know y such that SHA2(y)=z"	produces 2 polynomials: Q(X,Y,T,W), R(X) and	I know 4 polynomials of degree d - A(x), B(x), C(x), D(X) - such that:	If A, B, C, D do not satisfy <mark>THIS</mark> ,
2	degree bound d	Q(X, A(X), B(X+1), C(2*X))=D(X) * R(X)	then nearly all x expose Bob's lie

1) Arithmetization

Assuming Theorem, we get a scalable proof system for Bob's original claim:

- 1. Apply reduction, ask Bob to provide access to A,B,C,D of degree-d
- 2. Sample random x and accept Bob's claim iff equality holds for this x

Pre-arithmetization claim	Reduction	Post-arithmetization claim	Theorem
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2) Low degreeness

Assuming Theorem, we get a scalable proof system for Bob's original claim:

- 1. Apply reduction, ask Bob to provide access to A,B,C,D of degree-d
- 2. Sample random x and accept Bob's claim iff equality holds for this x

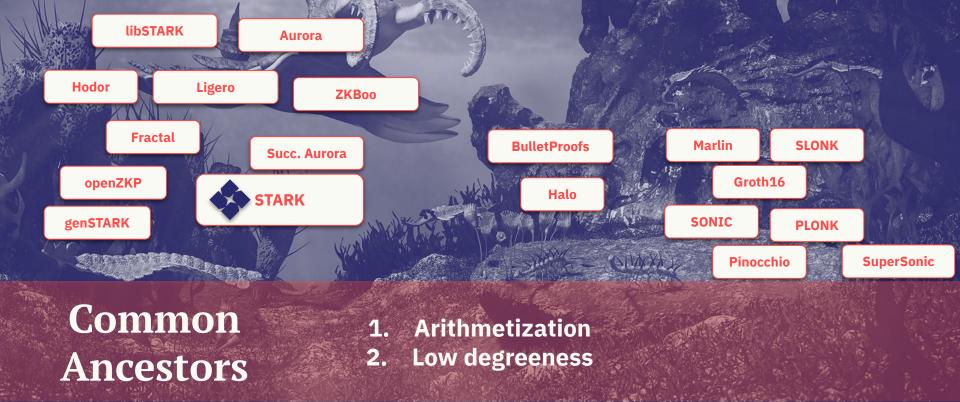
New Computational Integrity problem: Force Bob to answer all queries according to some quadruple of degree-d polynomials

Post-arithmetization claim	Theorem
I know 4 polynomials of degree d - A(x), B(x), C(x), D(X) - such that:	If A, B, C, D do not satisfy THIS,
Q(X, A(X), B(X+1), C(2*X))=D(X) * R(X)	then nearly all x expose Bob's lie

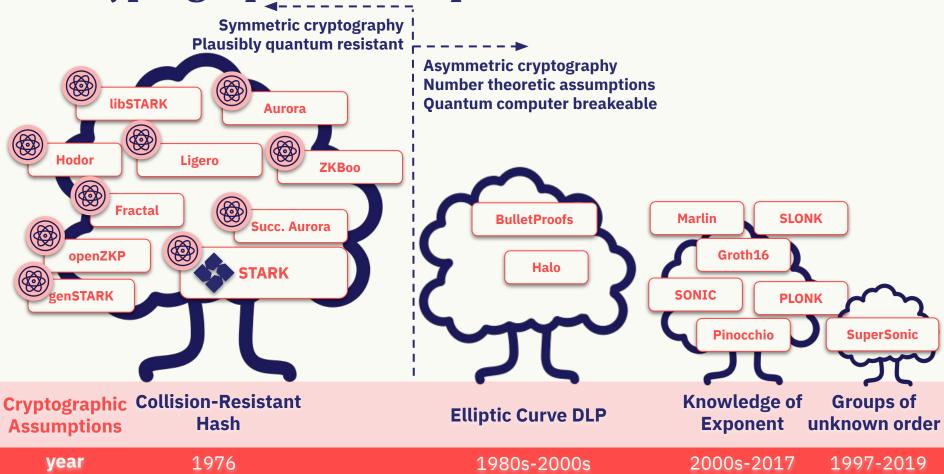


Differentiating Factors

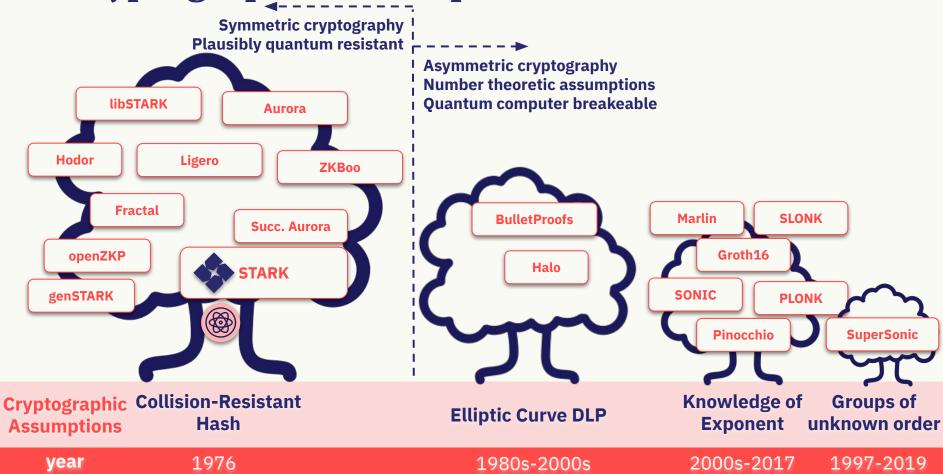
- **1.** Arithmetization Method
- 2. Low degreeness enforcement
- 3. Cryptographic assumptions used to get 2

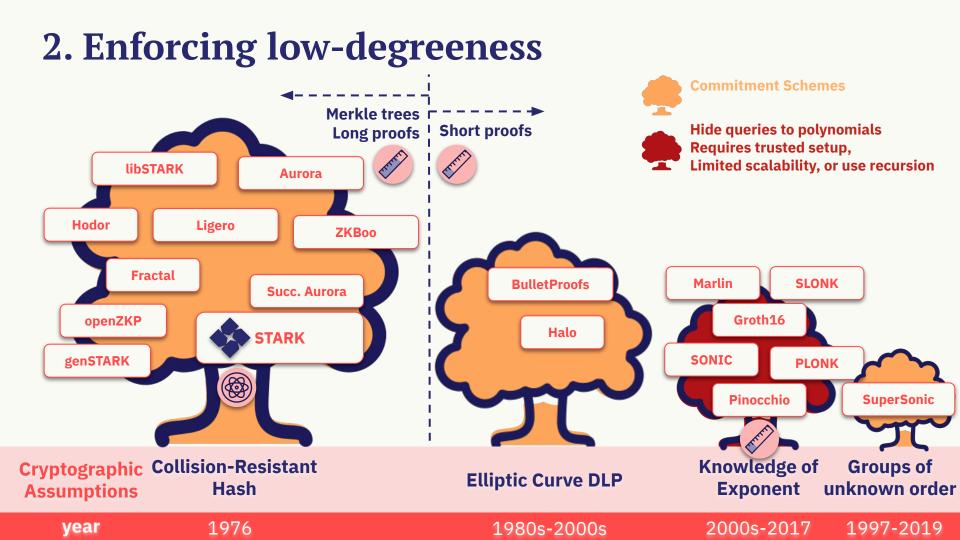


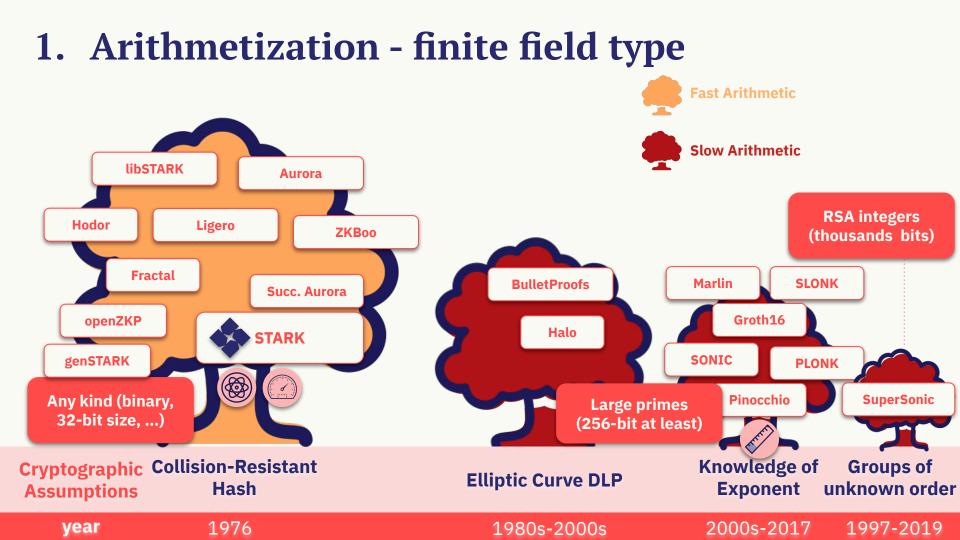
3. Cryptographic Assumptions

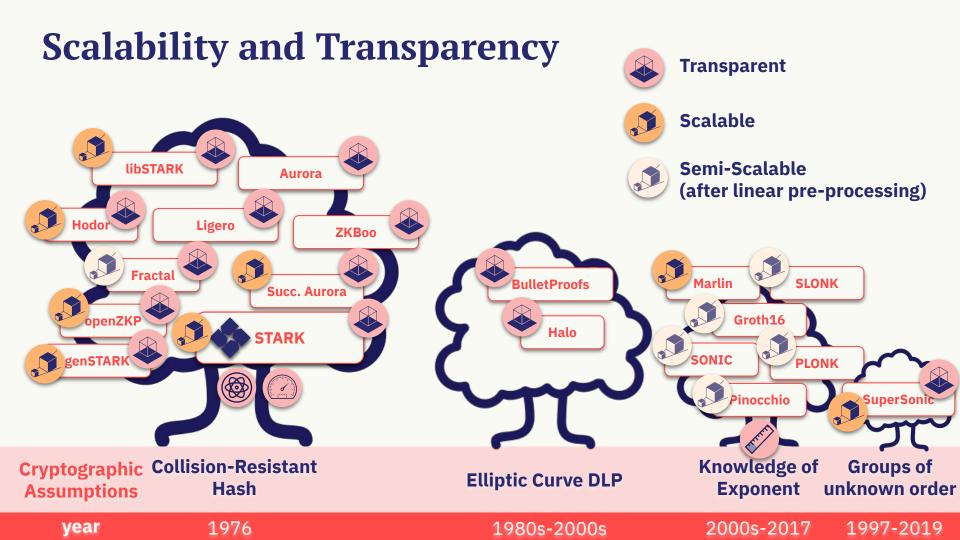


3. Cryptographic Assumptions



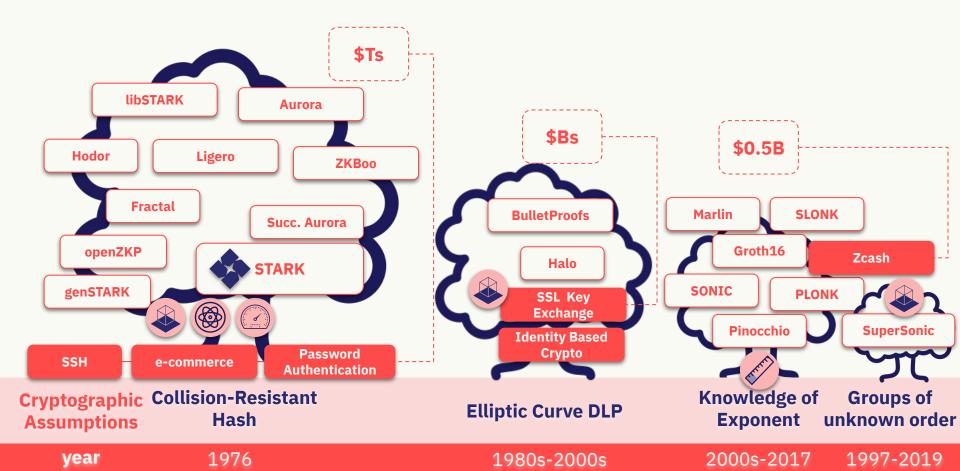




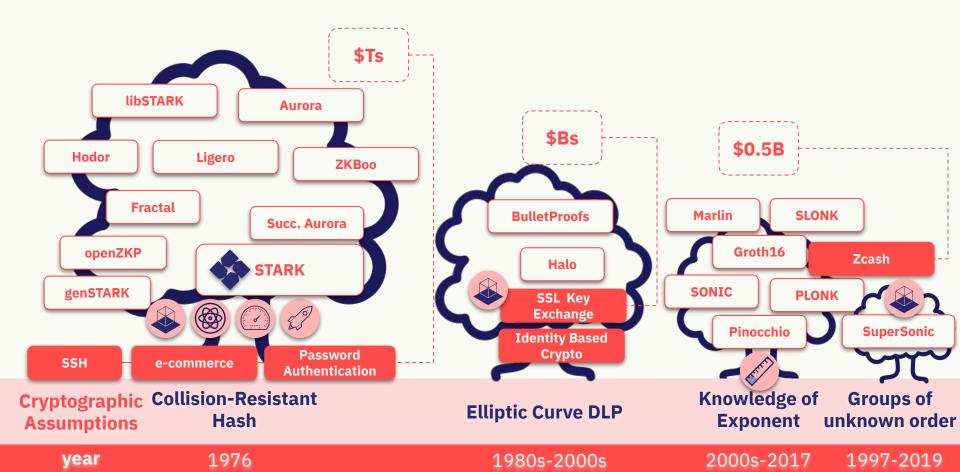


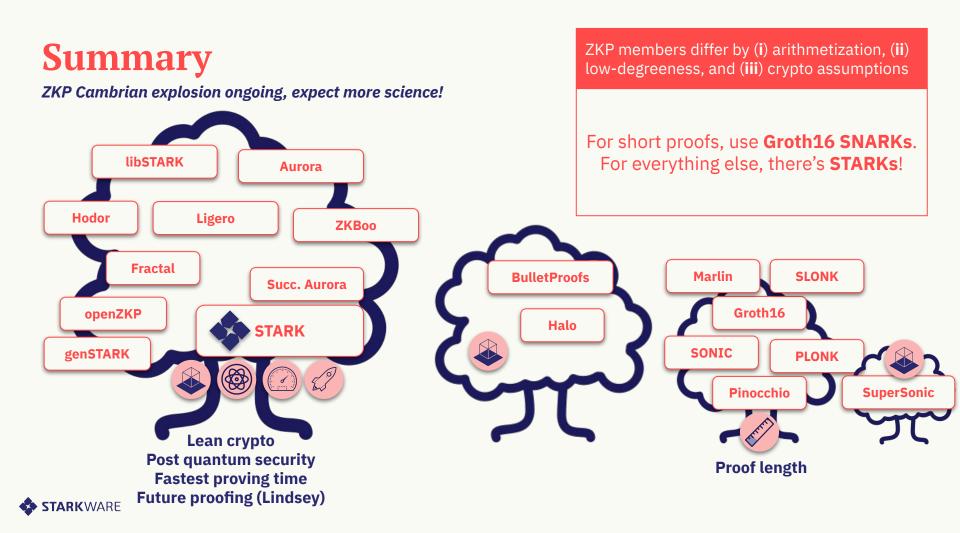
"The future life expectancy of some non-perishable things like a technology or an idea is proportional to their current age" libSTARK ~ The Lindy Effect / Nassim Taleb Aurora Hodor Ligero ZKBoo Fractal Marlin **SLONK BulletProofs** Succ. Aurora Groth16 openZKP Halo STARK SONIC genSTARK **PLONK Future-Proofing the Pinocchio SuperSonic** Financial Highway

ZKP Family Trees



ZKP Family Trees





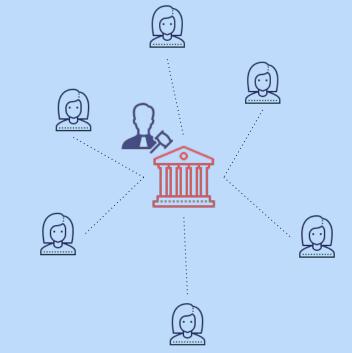


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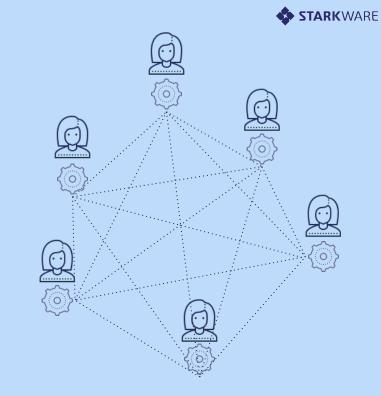
Trusted Party (e.g., Banks) = Delegated Accountability



Trust central party/auditor

Blockchains = Inclusive Accountability

Verify, Don't Trust

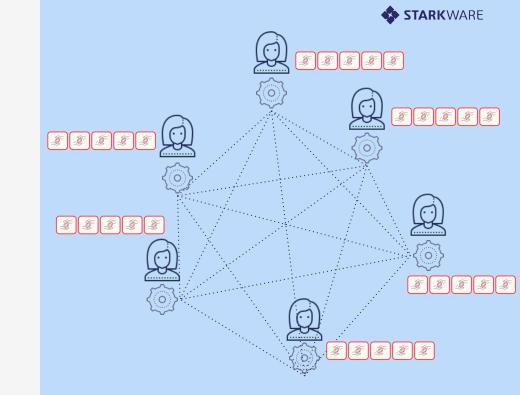


Verify (all transactions), don't trust

Blockchains = Inclusive Accountability

Sacrifice Privacy & Scalability

ZK-STARKs solve both problems



Verify (all transactions), don't trust

ZK-STARK Proofs

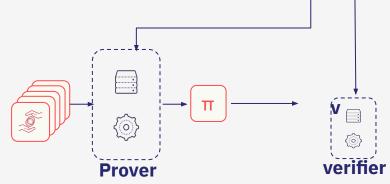


Privacy (Zero Knowledge, ZK) Prover's private inputs are shielded



Scalability

Exponentially small verifier running time* Nearly linear prover running time*



SALES REC	ЕІРТ	
1		
220 2017		
Description	Price	Amount
thach Salad	\$8.50	\$ 8.50
b Tagine e Rice	\$ 19.00	\$19.00
e Rice	\$ 4.00	\$ 16.00
ke	\$ 2.50	\$ 5.00
Ŷ	\$14.00	\$ 28.00
	Subtotal	\$76.50
	Tax	\$76.50

STARKWARE

VERIFIER Party checking proof (Customer)

Party producing proof

CI STATEMENT total=\$89.50

(Grocer)

Subtotal \$76.0 Tax: \$16.00 Tax: \$16.00 Total \$29.50 To

Date: May

Qty.

ZK-STARK Proofs

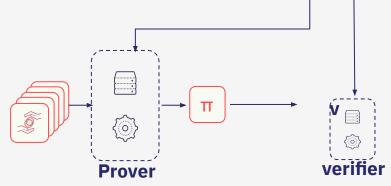


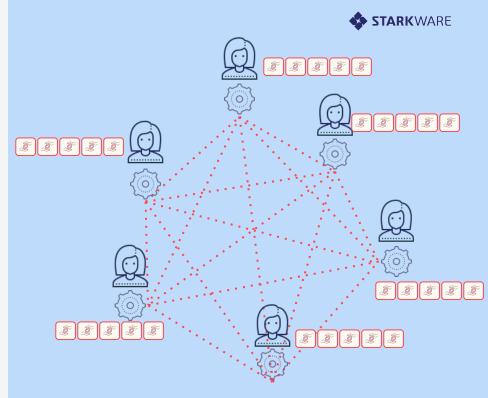
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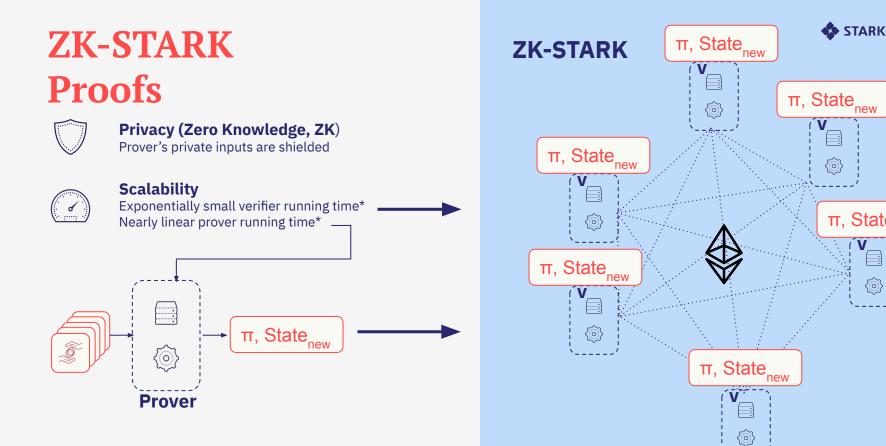
Scalability

Exponentially small verifier running time* Nearly linear prover running time*





Verify (all transactions), don't trust



Verify STARK proof, don't trust

STARKWARE

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Two L2 Offerings



Largest L2 by TPS

Roughly same rate as Ethereum, rising







StarkWare



Products

StarkEx Scalability Engine StarkNet STARK-Rollup







Pedigree

Invented ZK-STARK, FRI, Cairo, SHARP, Validium, Volition, ...

O(0)

70 Team members





Mission Bringing scalability & privacy to a blockchain near you



\$160M Funding (equity + EF grant)







Launched - June 2020

\$420	B	106M	>100K	
Cumulative T	rading	Tx Settled	Registered Users	
36M NFTs Minte		600K T Mints/Proof	486 Gas/tx	
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StarkNet

Decentralized Permissionless Validity-Rollup offering scalable & secure Ethereum-like state

L2

SMART CONTRACTS

GENERAL COMPUTATION

COMPOSABILITY



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- Learn
 - StarkNet/Cairo 101
 - Hello StarkNet!

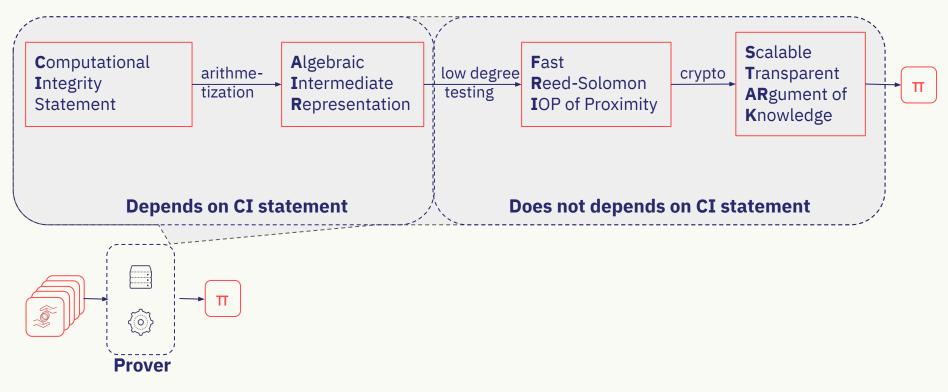
StarkNet Resources *

- Explore the Ecosystem
 - ➤ <u>StarkNet.io</u>
- Awesome StarkNet
- Stay up to date
 - ➤ StarkNet <u>roadmap</u>
 - StarkNet <u>unofficial newsletter</u>



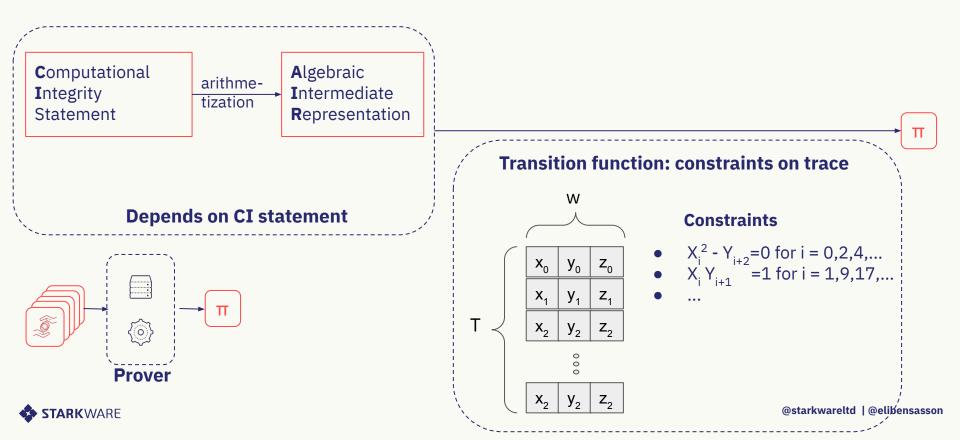
How to build a STARK?

How to build an AIR-FRI STARK





How to build an AIR-FRI STARK



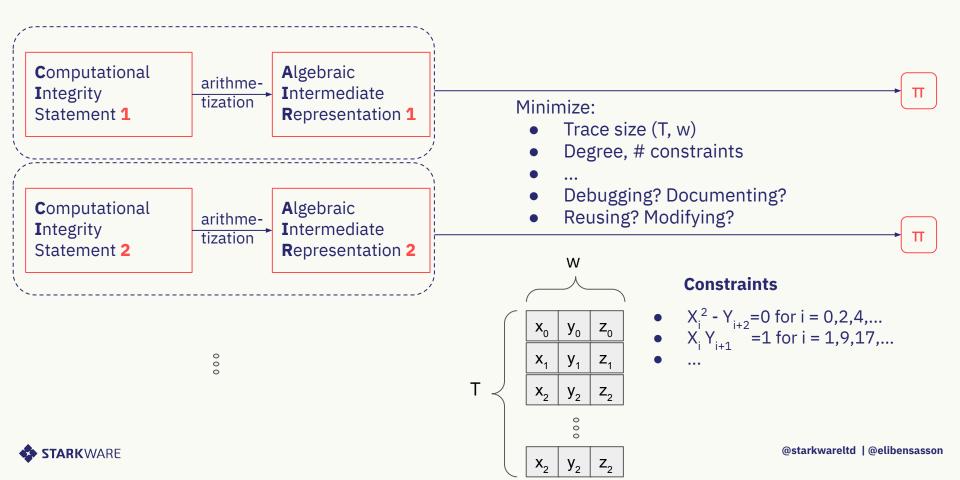
AIR Visualizer

STARKWARE

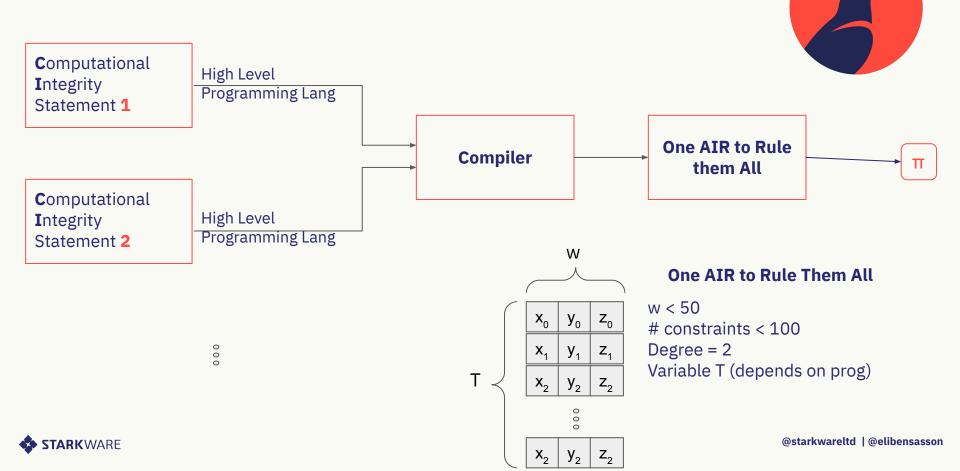
CO	C1	C2	C3	C4	C5	C6	C7	C8	C9	C10
IA	Step0_A	Step1_A	Step2_A	Step3_A	Step4_A	Step5_A	Step6_A	Step7_A	Step8_A	IB
Step9_A	Step0_A	Step1_A	Step2_A	Step3_A	Step4_A	Step5_A	Step6_A	Step7_A	Step8_A	Step9_B
ste A 9_A	Step0_A	Step1_A	Step2_A	Step3_A	Step4_A	Step5_A	Step6_A	Step7_A	Step8_A	Ste C9_B
Step9_A		ep0_a • (mat00 * (A - B	s) + mat01 * (C	- D)) * (mat00 *	(A - B) + mat0	1 * (C - D)) * (m	at00 * (A - B) +	mat01 * (C - D)) = 0 P8_A	Step9_B
Step9_A	Step0_A	Step1_A	Step2_A	Step3_A	Step4_A	Step5_A	Step6_A	Step7_A	Step8_A	Step9_B
Step9_A	Step0_A	Step1_A	Step2_A	Step3_A	Step4_A	Step5_A	Step6_A	Step7_A	Step8_A	Step9_B
Step9_A	Step0_A	Step1_A	Step2_A	Step3_A	Step4_A	Step5_A	Step6_A	Step7_A	Step8_A	Step9_B
Step9_A	Step0_A	Step1_A	Step2_A	Step3_A	Step4_A	Step5_A	Step6_A	Step7_A	Step8_A	Step9_B

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ASIC-like STARK

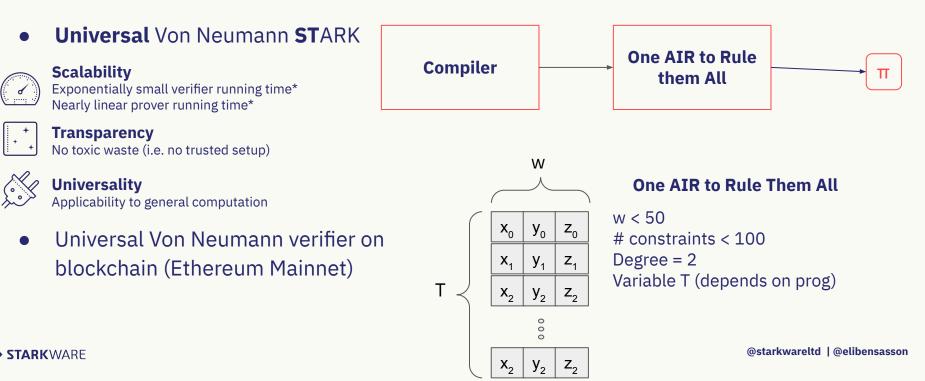


CPU AIR - CAIRo



Cairo Theory

Cairo is 1st





Cairo Theory

Cairo is 1st

• Universal Von Neumann STARK



Scalability Exponentially small verifier running time* Nearly linear prover running time*



Transparency No toxic waste (i.e. no trusted setup)



Universality Applicability to general computation

• Universal Von Neumann verifier on blockchain (Ethereum Mainnet)

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Resources:

Cairo landing page: https://cairo-lang.org/

Cairo whitepaper:

https://www.cairo-lang.org/cairo-whitepaper/

<u>Automated Theorem proving of Cairo</u> <u>soundness</u>: https://arxiv.org/abs/2109.14534





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 - a. <u>STARK 101 online course</u>: https://starkware.co/stark-101/
 - b. <u>STARK Math primer and whitepapers</u>: https://starkware.co/stark/



Questions?

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